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IP Payload Compression Protocol (IPComp)

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Abstract

This document describes a protocol intended to provide lossless compression for Internet Protocol datagrams in an Internet environment.

1. Introduction

IP payload compression is a protocol to reduce the size of IP datagrams. This protocol will increase the overall communication performance between a pair of communicating hosts/gateways ("nodes") by compressing the datagrams, provided the nodes have sufficient computation power, through either CPU capacity or a compression coprocessor, and the communication is over slow or congested links.

IP payload compression is especially useful when encryption is applied to IP datagrams. Encrypting the IP datagram causes the data to be random in nature, rendering compression at lower protocol layers (e.g., PPP Compression Control Protocol [RFC-1962]) ineffective. If both compression and encryption are required, compression MUST be applied before encryption.

This document defines the IP payload compression protocol (IPComp), the IPComp packet structure, the IPComp Association (IPCA), and several methods to negotiate the IPCA.

Other documents shall specify how a specific compression algorithm can be used with the IP payload compression protocol. Such algorithms are beyond the scope of this document.

1.1. Specification of Requirements

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC-2119].

2. Compression Process

The compression processing of IP datagrams has two phases: compressing of outbound IP datagrams ("compression") and decompressing of inbound datagrams ("decompression"). The compression processing MUST be lossless, ensuring that the IP datagram, after being compressed and decompressed, is identical to the original IP datagram.

Each IP datagram is compressed and decompressed by itself without any relation to other datagrams ("stateless compression"), as IP datagrams may arrive out of order or not arrive at all. Each compressed IP datagram encapsulates a single IP payload.

Processing of inbound IP datagrams MUST support both compressed and non-compressed IP datagrams, in order to meet the non-expansion policy requirements, as defined in section 2.2.

The compression of outbound IP datagrams MUST be done before any IP security processing, such as encryption and authentication, and before any fragmentation of the IP datagram. In addition, in IP version 6 [RFC-2460], the compression of outbound IP datagrams MUST be done before the addition of either a Hop-by-Hop Options header or a Routing Header, since both carry information that must be examined and processed by possibly every node along a packet’s delivery path, and therefore MUST be sent in the original form.

Similarly, the decompression of inbound IP datagrams MUST be done after the reassembly of the IP datagrams, and after the completion of all IP security processing, such as authentication and decryption.

2.1. Compressed Payload

The compression is applied to a single array of octets, which are contiguous in the IP datagram. This array of octets always ends at the last octet of the IP packet payload. Note: a contiguous array of octets in the IP datagram may be not contiguous in physical memory.

In IP version 4 [RFC-0791], the compression is applied to the upper layer protocol (ULP) payload of the IP datagram. No portion of the IP header or the IP header options is compressed.

In the IPv6 context, IPComp is viewed as an end-to-end payload, and MUST not apply to hop-by-hop, routing, and fragmentation extension headers. The compression is applied starting at the first IP Header Option field that does not carry information that must be examined and processed by nodes along a packet's delivery path, if such IP Header Option field exists, and continues to the ULP payload of the IP datagram.

The size of a compressed payload, generated by the compression algorithm, MUST be in whole octet units.

As defined in section 3, an IPComp header is inserted immediately preceding the compressed payload. The original IP header is modified to indicate the usage of the IPComp protocol and the reduced size of the IP datagram. The original content of the Next Header (IPv6) or protocol (IPv4) field is stored in the IPComp header.

The decompression is applied to a single contiguous array of octets in the IP datagram. The start of the array of octets immediately follows the IPComp header and ends at the last octet of the IP payload. If the decompression process is successfully completed, the IP header is modified to indicate the size of the decompressed IP datagram, and the original next header as stored in the IPComp header. The IPComp header is removed from the IP datagram and the decompressed payload immediately follows the IP header.

2.2. Non-Expansion Policy

If the total size of a compressed ULP payload and the IPComp header, as defined in section 3, is not smaller than the size of the original ULP payload, the IP datagram MUST be sent in the original non-compressed form. To clarify: If an IP datagram is sent non-compressed, no IPComp header is added to the datagram. This policy ensures saving the decompression processing cycles and avoiding incurring IP datagram fragmentation when the expanded datagram is larger than MTU.

Small IP datagrams are likely to expand as a result of compression. Therefore, a numeric threshold should be applied before compression, where IP datagrams of size smaller than the threshold are sent in the original form without attempting compression. The numeric threshold is implementation dependent.

An IP datagram with payload that has been previously compressed tends not to compress any further. The previously compressed payload may be the result of external processes, such as compression applied by an upper layer in the communication stack, or by an off-line compression utility. An adaptive algorithm should be implemented to avoid the performance hit. For example, if the compression of i consecutive IP datagrams of an IPCA fails, the next k IP datagrams are sent without attempting compression. If the next j datagrams are also failing to compress, the next k+n datagrams are sent without attempting compression. Once a datagram is compressed successfully, the normal process of IPComp restarts. Such an adaptive algorithm, including all the related thresholds, is implementation dependent.

During the processing of the payload, the compression algorithm MAY periodically apply a test to determine the compressibility of the processed data, similar to the requirements of [V42BIS]. The nature of the test is algorithm dependent. Once the compression algorithm detects that the data is non-compressible, the algorithm SHOULD stop processing the data, and the payload is sent in the original non-compressed form.

3. Compressed IP Datagram Header Structure

A compressed IP datagram is encapsulated by modifying the IP header and inserting an IPComp header immediately preceding the compressed payload. This section defines the IP header modifications both in IPv4 and IPv6, and the structure of the IPComp header.

3.1. IPv4 Header Modifications

The following IPv4 header fields are set before transmitting the compressed IP datagram:

Total Length

The length of the entire encapsulated IP datagram, including the IP header, the IPComp header and the compressed payload.

Protocol

The Protocol field is set to 108, IPComp Datagram, [RFC-1700].

Header Checksum

The Internet Header checksum [RFC-0791] of the IP header.

All other IPv4 header fields are kept unchanged, including any header options.

3.2. IPv6 Header Modifications

The following IPv6 header fields are set before transmitting the compressed IP datagram:

Payload Length

The length of the compressed IP payload.

Next Header

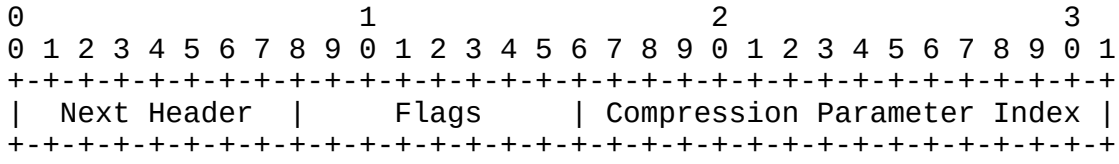
The Next Header field is set to 108, IPComp Datagram, [RFC-1700].

All other IPv6 header fields are kept unchanged, including any non-compressed header options.

The IPComp header is placed in an IPv6 packet using the same rules as the IPv6 Fragment Header. However if an IPv6 packet contains both an IPv6 Fragment Header and an IPComp header, the IPv6 Fragment Header MUST precede the IPComp header in the packet.

3.3. IPComp Header Structure

The four-octet header has the following structure:



Next Header

8-bit selector. Stores the IPv4 Protocol field or the IPv6 Next Header field of the original IP header.

Flags

8-bit field. Reserved for future use. MUST be set to zero. MUST be ignored by the receiving node.

Compression Parameter Index (CPI)

16-bit index. The CPI is stored in network order. The values 0-63 define well-known compression algorithms, which require no additional information, and are used for manual setup. The values themselves are identical to IPCOMP Transform identifiers as defined in [SECDOI]. Consult [SECDOI] for an initial set of defined values and for instructions on how to assign new values. The values 64-255 are reserved for future use. The values 256-61439 are negotiated between the two nodes in definition of an IPComp Association, as defined in section 4. Note: When negotiating one of the well-known algorithms, the nodes MAY select a CPI in the pre-defined range 0-63. The values 61440-65535 are for private use among mutually consenting parties. Both nodes participating can select a CPI value independently of each other and there is no relationships between the two separately chosen CPIs. The outbound IPComp header MUST use the CPI value chosen by the decompressing node. The CPI in combination with the destination IP address uniquely identifies the compression algorithm characteristics for the datagram.

4. IPComp Association (IPCA) Negotiation

To utilize the IPComp protocol, two nodes MUST first establish an IPComp Association (IPCA) between them. The IPCA includes all required information for the operation of IPComp, including the Compression Parameter Index (CPI), the mode of operation, the compression algorithm to be used, and any required parameter for the selected compression algorithm. The IPComp mode of operation is either a node-to-node policy where IPComp is applied to every IP packet between the nodes, or an ULP session based policy where only selected ULP sessions between the nodes are using IPComp. For each IPCA, a different compression algorithm may be negotiated in each direction, or only one direction may be compressed. The default is "no IPComp compression".

The IPCA is established by dynamic negotiations or by manual configuration. The dynamic negotiations SHOULD use the Internet Security Association and Key Management Protocol [ISAKMP], where IPSec is present. The dynamic negotiations MAY be implemented through a different protocol.

4.1. Use of ISAKMP

For IPComp in the context of IP Security, ISAKMP provides the necessary mechanisms to establish IPCA. IPComp Association is negotiated by the initiator using a Proposal Payload, which would

include one or more Transform Payloads. The Proposal Payload would specify a compression protocol in the protocol id field and each Transform Payload would contain the specific compression method(s) being offered to the responder.

In the Internet IP Security Domain of Interpretation (DOI), IPComp is negotiated as the Protocol ID PROTO_IPCOMP. The compression algorithm is negotiated as one of the defined IPCOMP Transform Identifiers.

4.2. Use of Non-ISAKMP Protocol

The dynamic negotiations MAY be implemented through a protocol other than ISAKMP. Such protocol is beyond the scope of this document.

4.3. Manual Configuration

Nodes may establish IPComp Associations using manual configuration. For this method, a limited number of Compression Parameters Indexes (CPIs) is designated to represent a list of specific compression methods.

5. Security Considerations

When IPComp is used in the context of IPSec, it is believed not to have an effect on the underlying security functionality provided by the IPSec protocol; i.e., the use of compression is not known to degrade or alter the nature of the underlying security architecture or the encryption technologies used to implement it.

When IPComp is used without IPSec, IP payload compression potentially reduces the security of the Internet, similar to the effects of IP encapsulation [RFC-2003]. For example, IPComp may make it difficult for border routers to filter datagrams based on header fields. In particular, the original value of the Protocol field in the IP header is not located in its normal positions within the datagram, and any transport layer header fields within the datagram, such as port numbers, are neither located in their normal positions within the datagram nor presented in their original values after compression. A filtering border router can filter the datagram only if it shares the IPComp Association used for the compression. To allow this sort of compression in environments in which all packets need to be filtered (or at least accounted for), a mechanism must be in place for the receiving node to securely communicate the IPComp Association to the border router. This might, more rarely, also apply to the IPComp Association used for outgoing datagrams.

6. References

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